

IT Compact Course

Hardware and Software

Internet and Web

Cryptography

Kaspar Etter, October 2011

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I. Hardware and Software

Hardware:

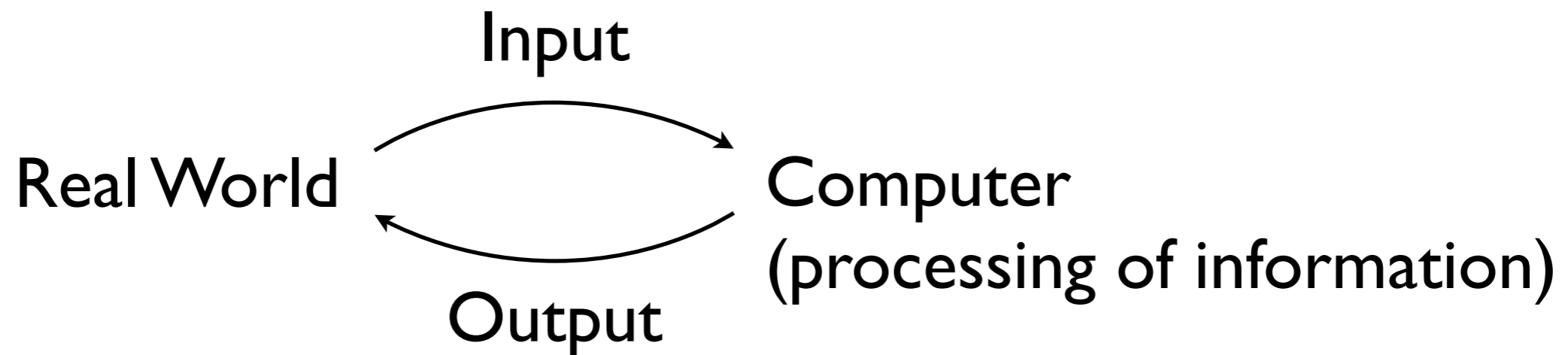
- Computers (physical)
- General purpose machines
- Expensive to design
- Expensive to copy
- Subjected to wear

Software:

- Programs (intangible)
- Special purpose instructions
- Expensive to design
- Free to copy
- Wear-free

I. Hardware and Software

Input/Output (I/O)



«Computer Science is no more about computers than astronomy is about telescopes.»

Edsger Dijkstra (misattributed)

I. Hardware and Software

Computers

The good news are that ...

- they do exactly what you tell them to do
- they do it very fast

The bad news are that ...

- they do exactly what you tell them to do
- they do it very fast

«To err is human, but to really mess things up you need a computer!»

I. Hardware and Software

Outline

I.1. Processor

I.2. Memory

I.3. Program

I.4. Operating System

I.5. Data Structures

I.6. Algorithms

} Hardware

} Software

} Programming

General reference and source for
further reading: www.wikipedia.org

I. Hardware and Software

I.1. Processor

- Central Processing Unit (CPU)
- Sequential processing of arithmetic and logical operations
- Data stored as binary numbers due to easy implementation in digital electronic circuitry using logic gates
- Only integers considered here (no floating-point numbers)
- A *digital* system uses *discrete* values, an *analog* system uses *continuous* values to represent information
(*Digital* comes from the Latin word *digitus*, meaning finger)
- A *bit* (a contraction of *binary digit*) is the basic unit of information in computing and is usually denoted as 0 and 1

1.1. Processor

Bits and Bytes

- 8 bits (b) = 1 byte/octet (B), allows to represent 256 values
- Unsigned 8-bit integer: 0 to 255; signed integer: -128 to 127
- Byte was the # of bits to encode a single character of text:
Basic addressable element in many computer architectures
- Processors manipulate bits in fix-sized groups named *words*

- Prefixes:

Decimal (SI)		Binary	
kilo (k)	10^3	kibi (Ki)	$2^{10} \approx 1.02 \cdot 10^3$
mega (M)	10^6	mebi (Mi)	$2^{20} \approx 1.05 \cdot 10^6$
giga (G)	10^9	gibi (Gi)	$2^{30} \approx 1.07 \cdot 10^9$
tera (T)	10^{12}	tebi (Ti)	$2^{40} \approx 1.10 \cdot 10^{12}$

Pointers, Registers and Flags

- Von Neumann architecture: Data & code in same memory
- Material based on Intel's instruction set architecture x86-32
Heavily simplified (no segmentation, addressing modes, etc.)
- Word length of 32 bits (= 4 bytes), from 0 to 4'294'967'295
- Instruction pointer: Memory address of next instruction
- 8 registers hold the current operands (the first 4 being general-purpose): EAX, EBX, ECX, EDX, ESI, EDI, ESP, EBP
- Carry, overflow and zero flag: Bits set after every operation
- Much of the trouble comes from backward compatibility

I.I. Processor

Operations

- Conceptually, there are only three types of operations:
 - Load data from memory to registers and store data from registers back to memory
 - Perform arithmetic and logical operations on registers
 - Control program flow with (conditional) jumps in code
- Memory is accessed with pointers to the desired locations
- Jumps can be absolute or relative in terms of memory address and often depend on the last executed operation
- The x86 instruction set comprises hundreds of operations

1.1. Processor

Assembly

- A low-level programming language that represents binary machine code in a human-readable form (with mnemonics)
- Needs to be translated into machine code for execution
- Example: Add together all numbers from 1 to 100

```
operator  operands  comment
  mov #0, sum ; set sum to 0
  mov #1, num ; set num to 1
loop: add num, sum ; add num to sum
label add #1, num ; add 1 to num
      cmp num, #100 ; compare num to 100
      ble loop ; if num <= 100, go back to 'loop'
      halt ; end of program. stop running
```

branch less or equal source and destination (register)

I.I. Processor

Pipelining

- Increase the instruction throughput by splitting the processing of an instruction into a series of independent steps (which increases the time to execute a single instruction)
- Issue instructions at the processing rate of the slowest step
- Maintain semantics for interdependent instructions and branches (branch prediction and speculative execution)

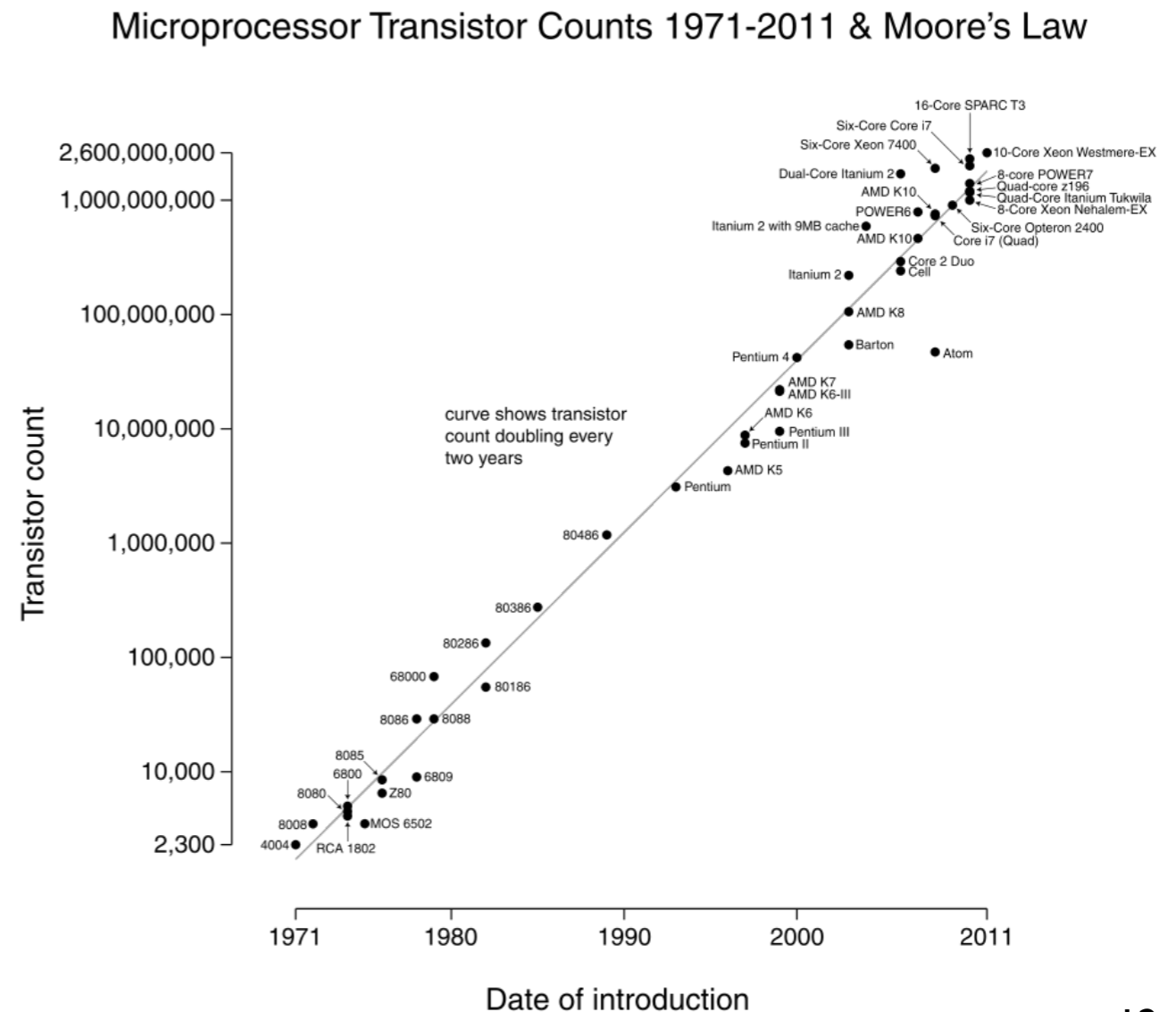
Instr No.	Pipeline Stage						
1	IF	ID	EX	MEM	WB		
2		IF	ID	EX	MEM	WB	
3			IF	ID	EX	MEM	WB
4				IF	ID	EX	MEM
5					IF	ID	EX
Clock Cycle	1	2	3	4	5	6	7

1. Instruction fetch
2. Instruction decode and register fetch
3. Execute
4. Memory access
5. Register write back

I.I. Processor

Moore's Law

- Trend described by Intel co-founder Gordon Moore, 1965:
The number of transistors on chips doubles every 2 years.
- Originally an observation and forecast, now a self-fulfilling prophecy
- Wirth's law, 1995:
Software is getting slower more rapidly than hardware becomes faster.



I. Hardware and Software

I.2. Memory

- A list of cells into which numbers can be placed or read
- The cells are numbered and can be addressed accordingly
- Hardware does not know the semantics of these numbers
- Memory is the bottleneck (limiting component of a system)

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Max-Bill-Platz
Richtung
Triemlispital

Gültig ab 12.12.2010

h	Montag-Freitag	Samstag	Sonn- und Feiertag	h
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5	09, 18, 24, 31, 39, 46, 53	09, 16, 24, 39, 54	16, 39, 54	5
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24	07, 17, 27, 37, 47, 57	09, 16, 26, 37, 53	09, 24, 37, 53	24

Schedule of bus and tram departures

1.2. Memory

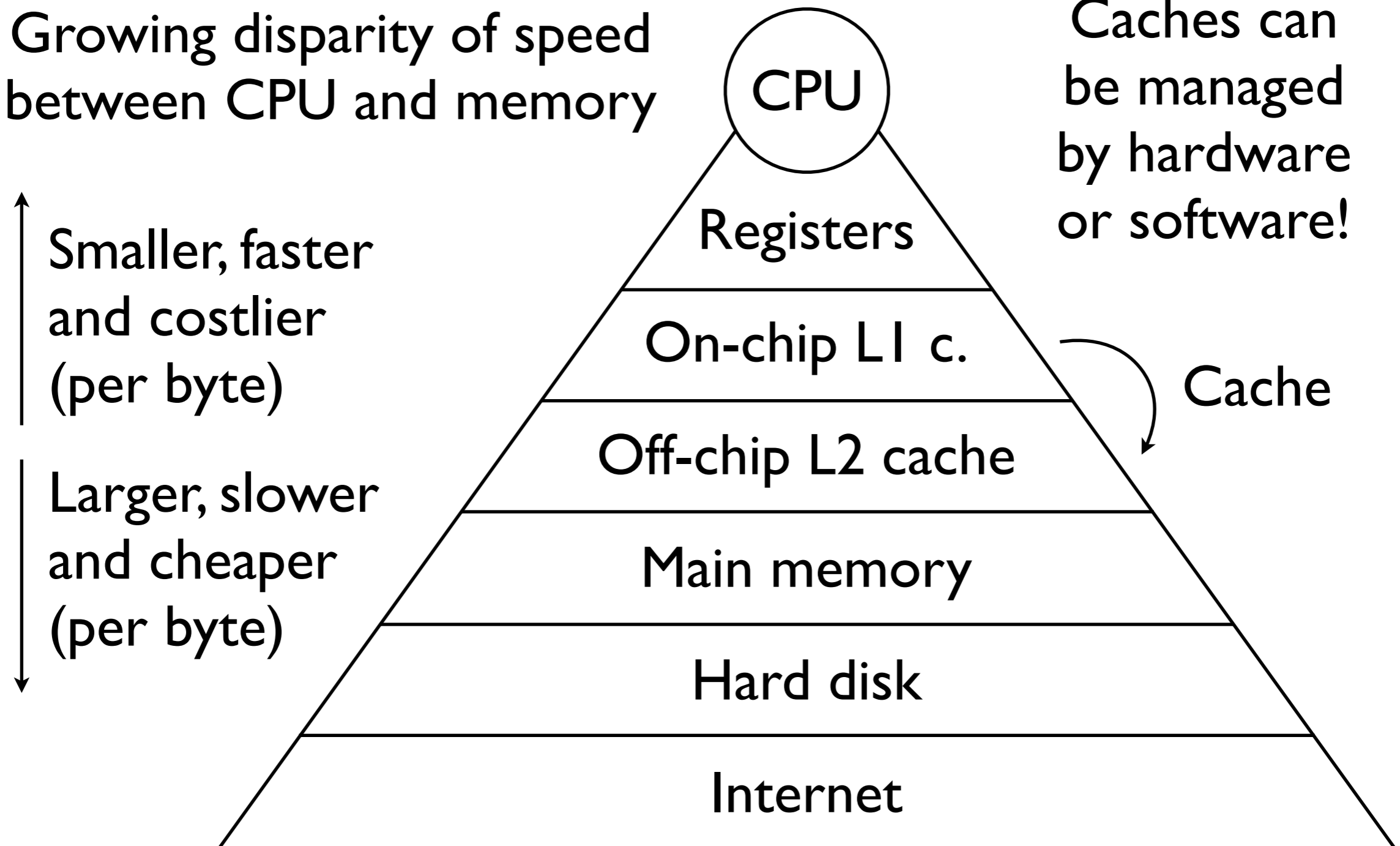
Caching

- A cache stores data for faster access in the future
- Cache hit: Requested data is contained in the cache
- Cache miss: Data has to be fetched from another location; Replaces a cache entry selected by the replacement policy
- The writing of data can be handled in two ways:
 - Write-through: Write to cache and memory concurrently
 - Write-back: Store the dirty cache entry on replacementEntries become stale, if somebody updates the original data
- Useful due to temporal and spatial locality of references
- The hit ratio ($\#hits / \#misses$) determines the performance

1.2. Memory

Memory Hierarchy

Growing disparity of speed between CPU and memory



1.2. Memory

Characteristics and Access

- ROM: Read-only memory, e.g. Compact Discs (CD-ROM); Used for firmware (low-level, hardware-specific software)
- RAM: Random-access memory, data can be accessed in any order (unlike disks); Often volatile storage (power supply!)
- Main memory connected to the CPU via a memory bus
- The memory management unit (MMU) calculates the actual memory address
- Disks: Seek time + rotation



I. Hardware and Software

I.3. Program

- A sequence of instructions that perform a specified task
- In its simplest form, takes some input and generates output
- The main qualities of software (programs) are:
 - **Correctness:** functional behavior according to specification
 - **Performance:** fast execution, low memory consumption
 - **Maintainability:** easy modification after initial development
 - **Reusability:** simple adaptation to new purposes/products
 - **Usability:** learnability, efficiency, memorability, satisfaction
 - **Security:** confidentiality, integrity, availability (cryptography)

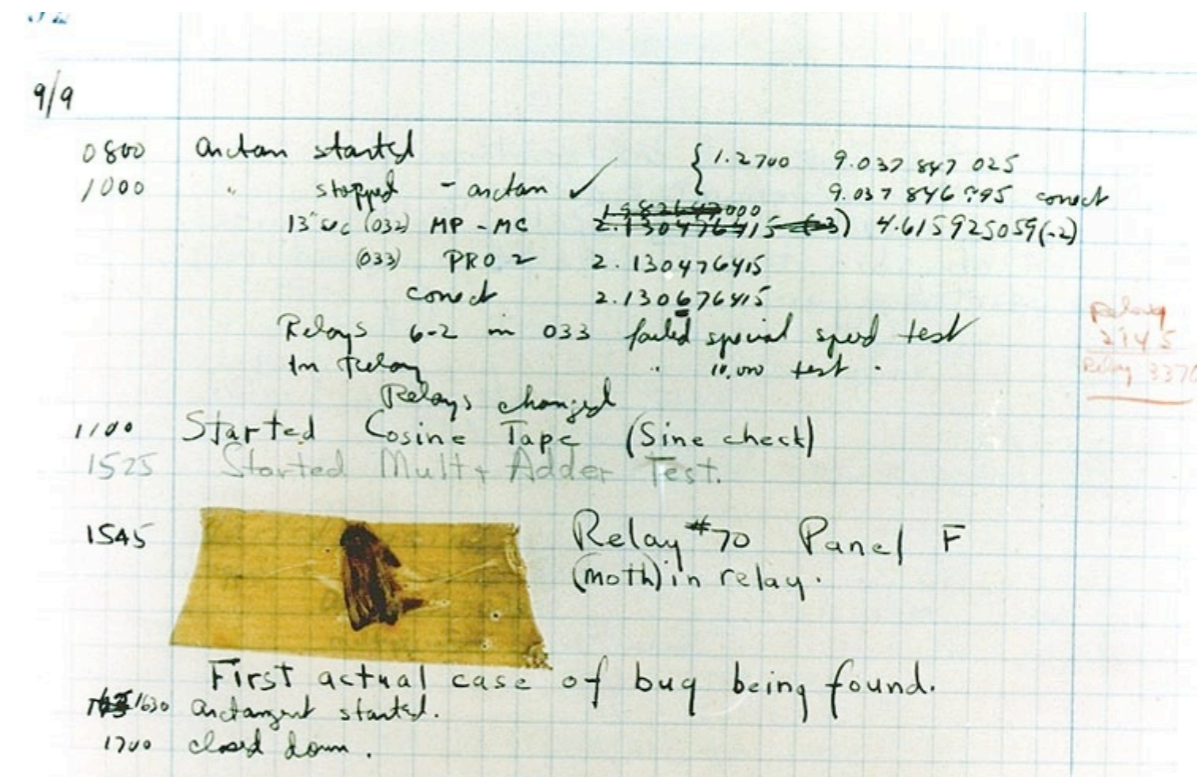
1.3. Program

Bugs

«Any feature is a bug unless it can be turned off.»

- Software errors cost the US economy \$60 billion annually, or about 0.6 percent of the gross domestic product (2002)
- US\$1 billion Ariane 5 rocket destroyed after takeoff (1996)
- Several patients were killed by the Therac-25 radiation therapy machine (1980s)
- Terminology: debug, buggy
- «Testing shows the presence, not the absence of bugs.»

Edsger Dijkstra



«First actual case of bug being found.» (1947)

1.3. Program

Programming Language

- An artificial language to give instructions to a computer
- Developing software in assembly language is error-prone
- Increase the programmer's productivity by providing a tool:
A program that mediates between man and machine
- Natural language is inappropriate: complex and ambiguous
- Formal languages consist of two parts: syntax and semantics
- Defined by a specification or a reference implementation
- Trade-off between abstraction/safety and expressiveness
- Benefit: High-level languages reduce platform dependency

1.3. Program

Source Code

- Text written in a high-level programming language that can be translated to binary machine code for execution
- Protected by copyright, protected form of free speech
- Example: “Hello world” in the C programming language

```
#include <stdio.h>

int main() {
    printf("Hello world!\n");
    return 0;
}
```

include standard I/O library

print formatted string

main method with return type integer
(called by the run-time environment)

code block

terminates statement

escape sequence for newline character

exit code indicating successful execution

1.3. Program

Compiler and Interpreter

- A compiler translates source code into machine code
- It checks the syntax and rejects invalid programs
- Semantic checks performed either statically (at compile time, e.g. types) or dynamically (at run time, e.g. arrays)
- Optimization: Constant propagation, common subexpression elimination, register allocation, instruction scheduling
- Interpreter: Execute code directly, present during execution
- Just-in-time compilation (JIT): Sections compiled 'on the fly'
- Reverse engineering: Decompilation (vs. code obfuscation)

1.3. Program

Text Terminal

- Command-line interface (CLI): Type commands to interact
- Graphical user interface (GUI): Manipulate visual elements
- Commands invoke programs with standard in- and output
- Interrupt their execution with ctrl-c (or ctrl-d on input)
- Structure: *command arguments* (syntax given by command)
- ‘>’ redirects output to a new file (‘>>’ appends to a file)
- ‘|’ chains output of left command & input of right command
- Prompt: Ready to accept commands, usually ends with ‘\$’
- Mac: Open ‘Applications/Utilities/Terminal’ and type ‘help’

I.3. Program

Demo: Integer Factorization

```
#include <stdio.h>
int main() {
    int number;
    printf("Number: ");
    scanf("%d", &number);
    while (number > 0) {
        printf("Factors: ");
        int factor = 2;
        while (factor * factor <= number) {
            if (number % factor == 0) {
                printf("%d, ", factor);
                number = number / factor;
            } else {
                factor = factor + 1;
            }
        }
        printf("%d\n", number);
        printf("Number: ");
        scanf("%d", &number);
    }
    return 0;
}
```

// Comments:
// Variable declaration
// Print to standard output
// Read a number from input
// Loop while condition true
// Assign value to variable
// Check remainder (modulo)

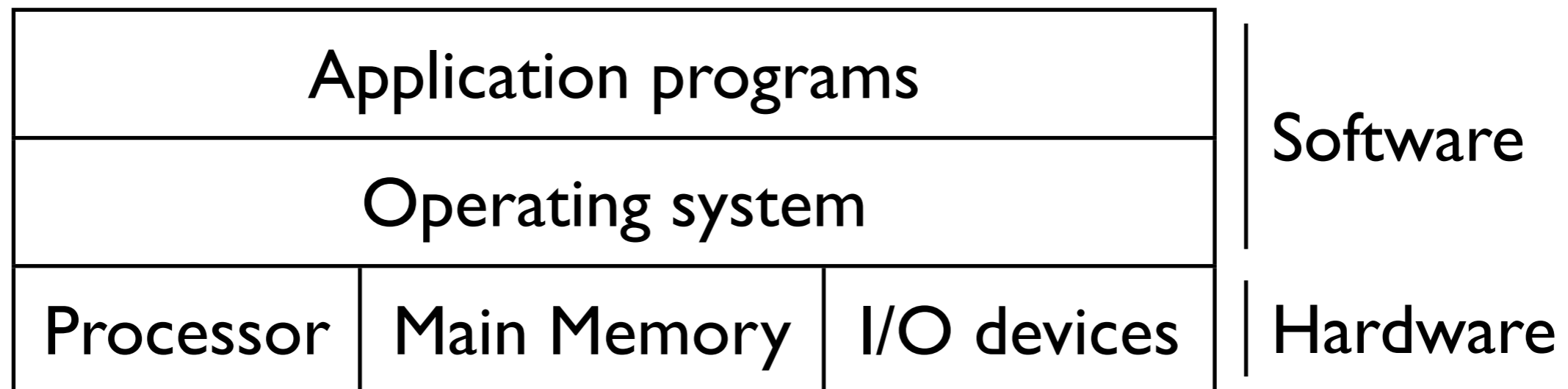
Store in a file 'code.c'
Compile with 'gcc code.c'
Run by typing './a.out'

Note: If you want to learn a language, learn Java and not C/C++!

I. Hardware and Software

I.4. Operating System

- The OS manages the hardware:
 - Abstraction: Simplify and standardize access to hardware (with so-called Application Programming Interface (API))
 - Duplication: Provide same resources to several programs
 - Protection: Ensure fairness and prevent misbehavior



I.4. Operating System

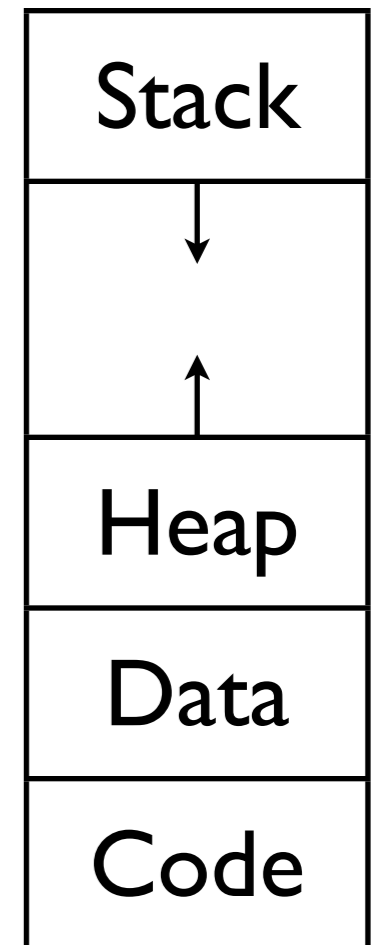
Process

- The OS creates for every running program a new process
- OS gives to each process the illusion of exclusive hardware: Execution without interruption, own memory and own I/O
- Processes can run concurrently by an interleaved execution
- Control transfer between processes with context switches
- Interrupts trigger special code in OS (privileged execution)
- Process scheduling done with regular hardware interrupts
- Processes have an owner and corresponding permissions
- A process can again have multiple execution units: Threads

I.4. Operating System

Virtual Memory

- Each process has its own virtual address space:
 - Code: Instructions copied from the executable
 - Data: Global variables (statically allocated)
 - Heap: Objects (dynamically allocated memory)
 - Stack: Tracks function calls and local variables
- Virtual memory split into blocks called pages
- OS allocates memory on demand at any location
- Every memory access gets translated into physical address
- If main memory is full, OS swaps inactive pages to hard disk



1.4. Operating System

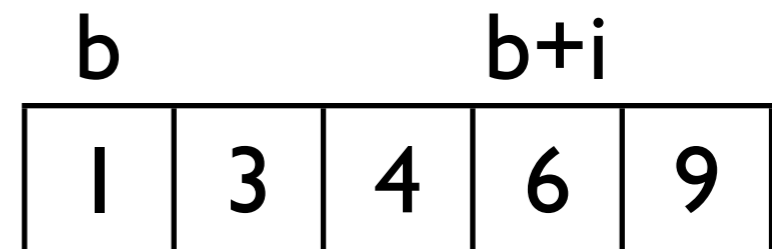
Partitioning

- Divide the hard disk drive into multiple logical storage units
- Booting is the process of loading the OS when starting up
- Multiple OSs can be installed on different partitions, loaded by the built-in basic input/output system (BIOS, “firmware”)
- File systems organize data to be retained after termination of a process: Store data permanently in units known as files
- Files consist of linked blocks, are structured by directories
- Specific FS layouts per partition, recover from corrupted FS
- Defragmentation reorganizes files into contiguous blocks

I. Hardware and Software

I.5. Data Structures

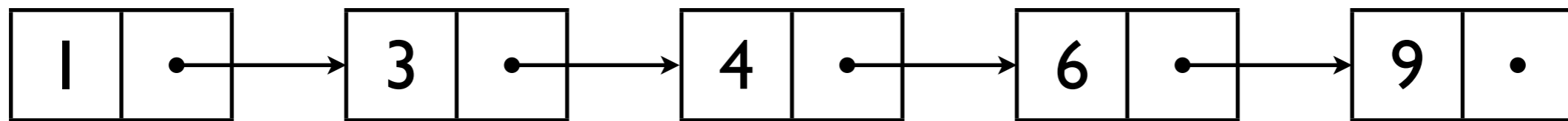
- A data structure is a way of storing and organizing data
- Different data structures suited to different applications
- Goal is time and space efficient manipulation of stored data
- Support the design of efficient algorithms (belong together)
- Arrays are a collection of elements in a continuous block
- Address of each element can be computed from its index
- Size fixed at allocation, insertion or removal needs copying
- Constant access, but linear insertion



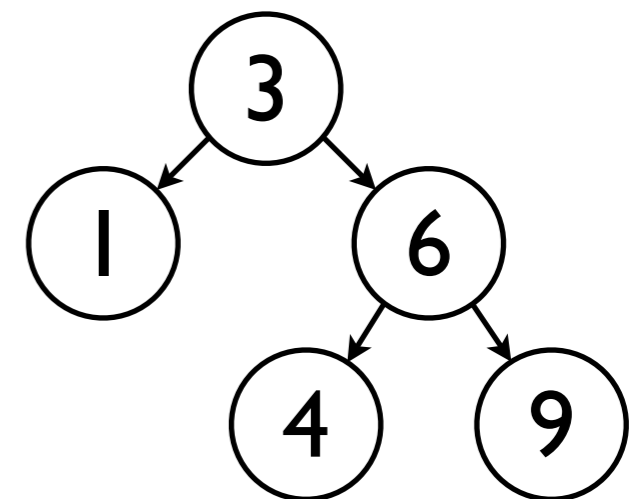
1.5. Data Structures

Lists and Trees

- A linked list is a sequence of elements stored in nodes
- Each node references the next node in the sequence
- Constant insertion and removal, but linear access time



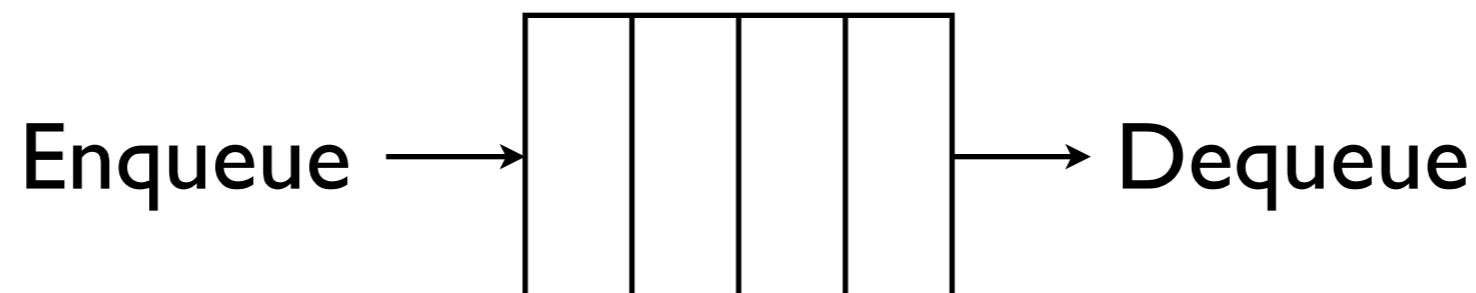
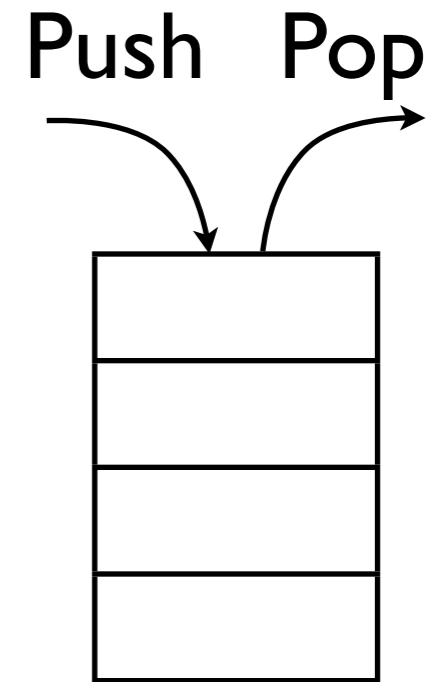
- A binary tree is a linked structure where every node has at most two child nodes
- Restriction that left nodes are smaller and right nodes are bigger allows binary search



1.5. Data Structures

Stacks and Queues

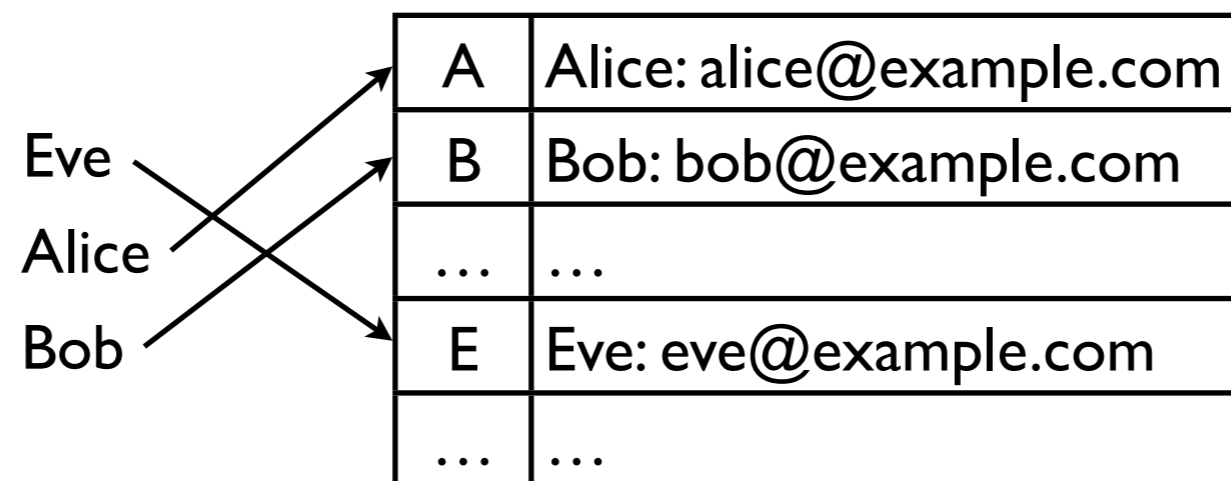
- A stack is a last in, first out (LIFO) collection
 - Only two operations provided: Push and pop
 - Implementation with an array or a linked list
 - Used to track open tasks (e.g. the call stack)
-
- A queue is a first in, first out (FIFO) collection (as a buffer)
 - Only two operations provided: Enqueue and dequeue



I.5. Data Structures

Hash Tables

- A hash table maps identifying values to associated values
- A hash function is used to transform the key into an index that indicates the corresponding value's position in an array
- A hash function maps values from a large to a small domain
- Collisions occur when different keys map to the same hash
- Widely used due to near-constant lookups (exc. collisions)



1.5. Data Structures

Buffer Overflow

- A buffer temporarily holds data that is moved from one place to another (implemented in hardware or software)
- Typically used for streaming (I/O) with variable rates
- Often implemented as a queue for simultaneous access
- Buffer overflows if incoming data exceeds storage capacity
- If not properly handled, adjacent memory gets overwritten
- Attacker overwrites return address or variable in call stack
- After decades of exploitation, still one of top vulnerabilities
- Solution: Use memory safe programming languages like Java

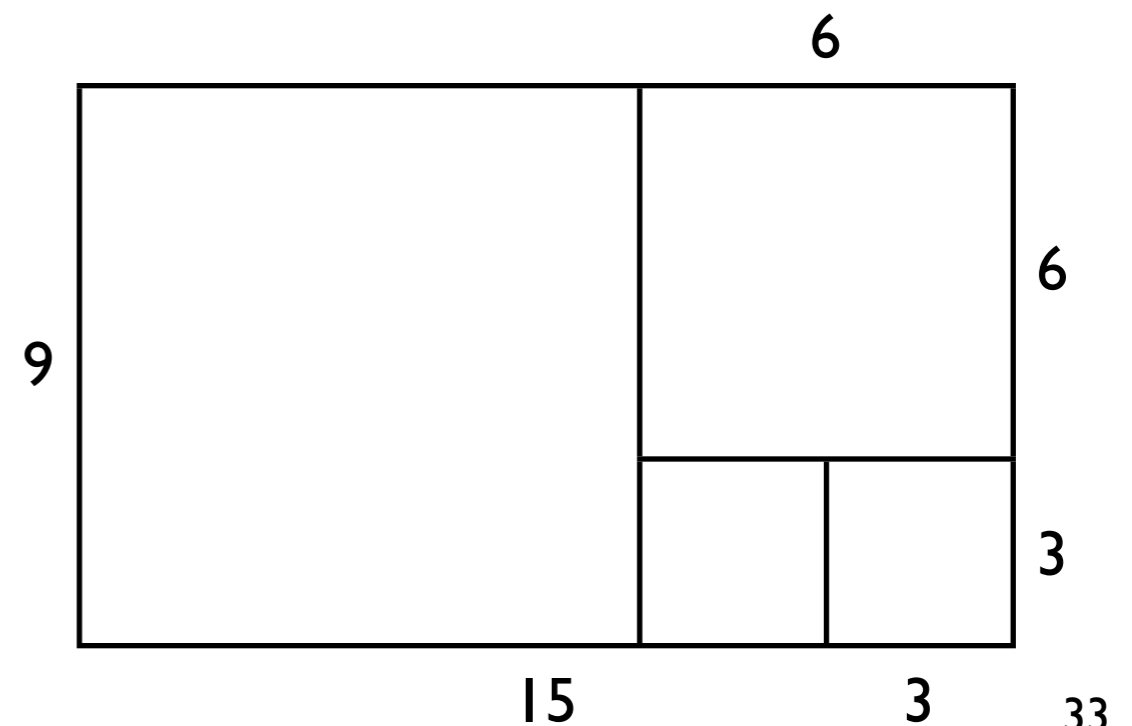
I. Hardware and Software

I.6. Algorithms

- An algorithm is a procedure for solving a specified problem in a finite number of steps (i.e. eventually producing output)
- Transitions between states do not have to be deterministic
- Brute-force: Naïve method of trying every possible solution
- Euclid's algorithm for the greatest common divisor (GCD):

```
function gcd(a, b)
  if a = 0
    return b
  while b ≠ 0
    if a > b
      a := a - b
    else
      b := b - a
  return a
```

a	b
15	9
6	9
6	3
3	3
3	0



1.6. Algorithms

Sorting

- Given a list of comparable objects, return them in order
- Bubble sort: Steps repeatedly through the list, compares adjacent items and swaps them if they are in wrong order
 - Visualizing the sorting, small elements bubble to the top
 - In-place algorithm: Only constant amount of extra storage
- Merge sort: Divides the list into halves, sorts them recursively and merges the results (divide & conquer algorithm)
 - Preserves the input order of equal elements (stable sort)
 - Requires a linear amount of additional storage space
- See: cs.usfca.edu/~galles/visualization/ComparisonSort.html

1.6. Algorithms

Time Complexity

- Estimating the processing time of algorithms, we are only interested in how they respond to changes in input size: Efficiency measured by how well they scale with input size
- Big O notation characterizes functions according to their growth rate by suppressing multiplicative constants and lower order terms (upper bound): e.g. $5n^3 + 3n$ is $O(n^3)$
- Constant time (or space) complexity expressed as $O(1)$
- $O(n)$ denotes a linear, $O(2^n)$ an exponential time algorithm
- Bubble sort is $O(n^2)$: n rounds of n comparisons and swaps
- Merge sort is $O(n \log n)$: $\log n$ rounds of linear merging

1.6. Algorithms

Complexity Theory

- Classifying problems according to their inherent difficulty; Determine the practical limits of what computers can do
- A complexity class is a set of problems of related difficulty:
 - P: Problems deterministically solvable in polynomial time
 - NP: Non-deterministically solvable in polynomial time, i.e. a solution to the problem can be verified in polynomial t.
- Clearly, $P \subseteq NP$, NP containing many important problems; Hardest problems in NP are NP-complete (reduction...)
- Example for NP-complete: The subset sum problem ($\sum = 0$)
- $P \stackrel{?}{=} NP$ problem: Efficient check implies efficient solution?

1.6. Algorithms

Computability Theory

- Asks what kind of problems can be solved algorithmically
- A decision problem is a question with a yes-or-no answer
- A decision problem solvable by an algorithm is decidable
- Halting Problem: Given a description of a program and a finite input, decide whether the program finishes running
- An algorithm is required to terminate (i.e. in finite time)
- Alan Turing proved in 1936 that no general algorithm to solve the halting problem for all possible pairs can exist
- Proof by reduction: New algo. would solve undec. problem

I. Hardware and Software

Concepts Learned Today

- Abstraction
- Algorithm
- Bootstrapping
- Bottleneck
- Caching
- Complexity
- Pipelining
- Specification
- Transparency

I. Hardware and Software

Clip of Today

My favorite comment:
Then again, there's the Brute Force
technique: Steal already sorted arrays
from third world countries. $O(1)$.

Barack Obama - Computer Science Question (1:25)



http://www.youtube.com/watch?v=k4RRi_ntQc8

Questions?